Projects Summer Term 2025

Supervisors: Alexander Dinges, Cass Alexandru, Michael Youssef, Sebastian Schloßer Kick-off: 30.04.2025

Organization

- Kick-off: Wednesday, 30.04.2025, 13:45 15:15, room 36-265
- Final presentations: presumably Friday, 25.07.2025, time and room TBA Mandatory attendance!
- Regular meetings with your supervisor

- Four topics to choose from
- One or two teams per topic (working independently)
- Four or five students per team (collaborating)

- We present you the topics
- You (physically) go to the supervisor whose topic you would like to work on
- There you form teams and exchange details with team members & supervisor

Data Management for ExClaim

Supervisor: Sebastian Schloßer

EXCLAIM

- Exercise management
 - $\circ \ \, {\sf Students}$
 - Groups (same exercise session)
 - Teams (collaborate on homework)
 - \circ Homework submission
 - $\circ~$ Homework grading
 - Automatic group assignment
- Exam management
- Automatic test execution

- Docker
- Consistent environment
- No risks for tutors
- Fast feedback (for students and as guidance while grading)
- Can execute test cases without providing the source code to students

- Softech Achievement Tracking System (STATS)
 - $\circ~$ Manage students, groups, teams, exams
 - No submission (on paper or via e-mail)
 - $\circ~$ No automatic group assignment
- Optimus: Automatic group assignment
- ExClaim & RTE: For submission & testing

Three separate systems!

- Users management only in STATS
- Linking data via student id

- Integrate all features in one system
- Optimize database layout
- Use code-generation for type-safe database access
- Simplify building the system
- New Frontend in development

- Initialize Database with dummy data
- Facilitates development
- For demonotration purposes
- (Cannot use real data!)

- Retrieve all data for a single course (for offline evaluation)
- Retrieve all data for a single user (to fulfill GDPR requests)
- Retrieve essential data for a single course (Overall homework points per student, exam assessment, but no uploads or single sheet points)
- Format: JSON or XML

- Delete entire course
- from database (recursively all dependent data) and stored files (homework uploads)
- Required to delete data after retention period
- To not lose exam admissions, archived export of essential data is important

Technologies to Learn / Know

- Git
- Backend
 - Java
 - \circ Spring Boot
 - $\circ\,$ jOOQ (Code Generator for type-safe database access)
- Frontend
 - $\circ~$ TypeScript / Javascript
 - $\circ \ \mathsf{Vue}$
 - $\circ~$ HTML, CSS

Search tool for Agda

Supervisor: Alexander Dinges

Your task

- CLI search tool for an Agda repository
- Find function signatures/propositions/types together with their names given a part of the type
- Reasonably fast

Minimum requirements for getting your credit points

- Given a signature $t_1 \rightarrow t_2 \rightarrow \ldots t_n$ find all (complete) signatures (in a given Agda code base) that contain t_1, \ldots, t_{n-1} as parameter type and t_n as return type.
 - ▶ t₁,..., t_n match up to variable renaming (consistently) and number representation
 - I.e. find 1 ≺ n → 0 ≺ n/2 when given (succ zero) ≺ a → zero ≺ a/2
- Don't find irrelevant stuff.
- Reasonable ranking and speed.

Minimum requirements for getting your credit points

Moreover, choose at least 2 of the following points:

- Make the CLI user-friendly: Agda input method, colors, search history, ...
- Allow more flexible search strings: Partial matching, underscore patterns, ...
- Good ranking.
- Allow more unification
- Your own idea?

Technologies to learn/know

- Git
- A tiny bit of Agda
- You can use (almost) any programming language you like

Developing a web app using Haskell

Supervisor: Michael Youssef

Why Haskell?

- Haskell is a pure language
- No unintended side effects
- ADTs
- Modularity
- Laziness
- Referential transparency

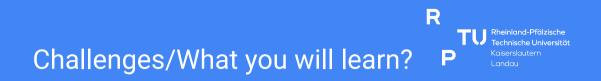


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Why Haskell?

- Haskell is a pure language
- No unintended side effects
- ADTs
- Modularity
- Laziness
- Referential transparency

- No null-pointer exceptions
- No IO exceptions
- Elegant approach towards handling IO
- No unhandled cases



- Type classes
- Practical usages of functors, applicatives and monads
- Monad transformers
- Type families
- Template Haskell
- Debugging with lazy evaluation

Challenges/What you will learn? P Landau

- Type classes
- Practical usages of functors, applicatives and monads
- Monad transformers
- Type families
- Template Haskell
- Debugging with lazy evaluation

- Network IO
- Database storage
- Caching
- Writing a fully fledged application in Haskell

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Why should you care?

- ECTS....
- Learn some practical applications of the stuff you learned in FP
- Knowledge you gain is transferable to other programming language paradigms

LATEX-formatted Execution Traces of Algorithms on Automata & Grammars

Supervisor: Cass Alexandru

LATEX-formatted Execution Traces of Algorithms on Automata & Grammars

Problem Statement

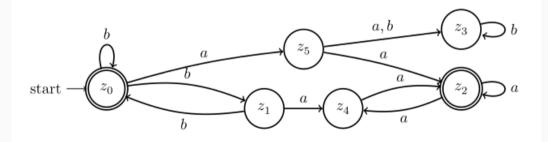
- Lecture "Formal Languages and Computability"
- $\bullet\,$ Students learn to execute several algorithms on automata & grammars by hand
- I (the client) want a software tool that:
 - $\circ~$ Given an input object (an automaton or grammar)
 - Executes one of a number of algorithms on it and outputs its execution trace (showing its steps) in the format used in the lecture
- This tool would allow me to more quickly iterate exercise ideas & easily and confidently generate correct reference solutions without time-consuming and error-prone manual calculation

• Automata:

- \circ Minimization
- Determinization (Powerset/Rabin-Scott Construction)
- Product Automaton
- Execution trace ((nondeterministic) stack automaton)
- Execution trace (Turing machine)
- Grammars:
 - CNF (Chomsky Normal Form) algorithm for context-free grammar
 - CYK Algorithm for bottom-up parsing of words in a cf grammar
 - Derivation of word from starting symbol in case of language membership

P2-1 Läufe und Potenzmengenkonstruktion

Für diese Aufgabe betrachten wir folgenden NFA $\mathcal C$ über dem Alphabet $\{a,b\}$

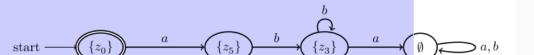


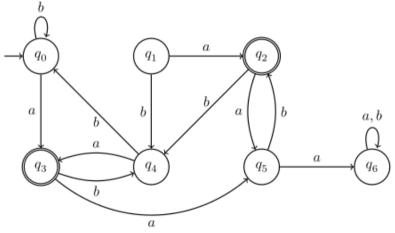
LÖSUNGSVORSCHLAG:

Potenzmengenkonstruktion, wobei nur die erreichbaren Zustände betrachtet werden:

Zustand/Nachfolge bei	a	b
$\{z_0\}$	$\{z_5\}$	$\{z_0, z_1\}$
$\{z_5\}$	$\{z_2, z_3\}$	$\{z_3\}$
$\{z_0, z_1\}$	$\{z_5, z_4\}$	$\{z_0, z_1\}$
$\{z_2, z_3\}$	$\{z_2, z_4\}$	$\{z_3\}$
$\{z_3\}$	Ø	$\{z_3\}$
$\{z_5, z_4\}$	$\{z_2, z_3\}$	$\{z_3\}$
Ø	Ø	Ø
$\{z_2, z_4\}$	$\{z_2,z_4\}$	Ø

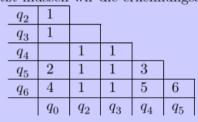
Startzustand $\{z_0\}$ und Endzustände: $\{z_0\}, \{z_0, z_1\}, \{z_2, z_3\}, \{z_2, z_4\}$





LÖSUNGSVORSCHLAG:

Wir müssen erst einmal jene Zustände entfernen, die aus dem Startzustand nicht erreichbar sind. In dem Fall ist das $\{q_1\}$. Jetzt müssen wir die erkennungsäquivalenten Zustände finden:

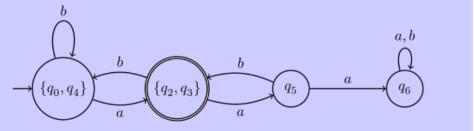


- (1) Markiere Paare von Endzuständen und Nicht-Endzuständen
- (2) Markiere $\{q_5,q_0\}$ we
gen $\delta(q_5,b)=q_2, \delta(q_0,b)=q_0$ und $\{q_2,q_0\}$ markiert
- (3) Markiere $\{q_5,q_4\}$ we
gen $\delta(q_5,b)=q_2, \delta(q_4,b)=q_0$ und $\{q_2,q_0\}$ markiert
- (4) Markiere $\{q_6,q_0\}$ we
gen $\delta(q_6,a)=q_6, \delta(q_0,a)=q_3$ und $\{q_6,q_3\}$ markiert
- (5) Markiere $\{q_6,q_4\}$ we
gen $\delta(q_6,a)=q_6, \delta(q_4,a)=q_3$ und $\{q_6,q_3\}$ markiert
- (6) Markiero [a, a] we go $\delta(a, b) = a$ $\delta(a, b) = a$ und [a, a] markiert

Wir können also die folgenden Zustände verschmelzen:

- *q*₀ mit *q*₄
- q_2 mit q_3

Und erhalten folgenden Minimalautomat:



P7-1 Chomsky Normalform (CNF) Wandeln Sie die Grammatik $\mathcal{G}_1 = (\{S, A, B\}, \{a, b\}, P, S) \text{ mit } P:$ $S \to BA \mid B \mid ABBA$ $A \to aa \mid \varepsilon$

$$B \rightarrow BA \mid bb$$

in Chomsky Normalform (CNF) um. Verwenden Sie dabei das Verfahren aus der Vorlesung.

LÖSUNGSVORSCHLAG:

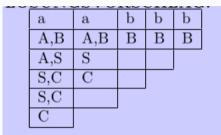
Wir befolgen die 4 Schritte, die in der Vorlesung gezeigt worden sind:

1. Eliminieren von ε Produktionen:

$$\begin{split} S &\to BA \mid B \mid BB \mid ABB \mid BBA \mid ABBA \\ A &\to aa \\ B &\to BA \mid B \mid bb \end{split}$$

$$\begin{split} S &\to AB \mid AC \\ A &\to AA \mid a \\ C &\to SB \\ B &\to a \mid b \end{split}$$

Benutzen Sie den CYK Algorithmus, um zu testen, ob *aabbb* und *aaabbb* $\in L(\mathcal{G}_2)$. Falls das Wort in der Sprache enthalten ist, geben Sie eine Ableitung an. Sie können folgende Tabellen benutzen:



 $aabbb \notin L(\mathcal{G}_2),$ weil das StartsymbolSnicht in dem untersten Kästchen vorkommt.

LÖSUNGSVORSCHLAG: PDA

 $\mathcal{P}_1 = (\{q_0, q_1, q_2, q_3\}, \{a, b, c, d\}, \{\#, A, B\}, q_0, \Delta, \#)$ mit $\Delta:$

$$\begin{split} (q_0,a,\#) &\rightarrow (q_0,A\#) \\ (q_0,a,A) &\rightarrow (q_0,AA) \\ (q_0,b,A) &\rightarrow (q_1,BA) \\ (q_1,b,B) &\rightarrow (q_1,BB) \\ (q_1,c,B) &\rightarrow (q_2,\varepsilon) \\ (q_2,c,B) &\rightarrow (q_2,\varepsilon) \\ (q_2,d,A) &\rightarrow (q_3,\varepsilon) \\ (q_3,d,A) &\rightarrow (q_3,\varepsilon) \\ (q_3,\varepsilon,\#) &\rightarrow (q_3,\varepsilon) \end{split}$$

Wörter in L_1 :

LATEX-formatted Execution Traces of Algorithms on Automata & Grammars

Execution

- At least one member should have experience with the following:
 - B2 level German (reference solns are generated in German, though you won't need to write much text yourself, I provide templates)
 - $\circ~\mbox{\sc BT}_{\mbox{\sc EX}}$ (again, I mostly provide templates and will also provide guidance)
 - $\circ~$ The algorithms in question from the "Formal Languages" lecture

- Regular meetings with me to keep shared understanding of requirements synched
- There will be milestones for deliverables with intermediate functionality
- Development should use Gitlab Issues, Issue Boards and, ideally, CI/CD, with provisions made for shared, replicable build and dev environments across the team and CI

- 23.05: DFA Minimization
- 30.05: CNF
- 06.06: CYK
- 13.06: PDA trace
- 20.06: TM trace
- 27.06: NFA Determinization
- 04.07: Product Automaton

Divide up responsibilities & pipeline tasks!